

Symbolic Computation in Automated Program Reasoning

Laura Kovács

for(synte,   Informatics

Automating Program Reasoning



Automating Program Reasoning

(ex. ~250kLoC, Vampire prover)



Automating Program Reasoning

```
a=0, b=0, c=0;  
while (a<n) do  
  if A[a]>0 then B[b]=A[a]+h(b); b=b+1;  
  else C[c]=A[a]; c=c+1;  
  
  a=a+1;  
end do
```

Automating Program Reasoning

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    else C[c]=A[a]; c=c+1;  
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end do
```

Program property:

$(\forall p)(0 \leq p < b \Rightarrow$

$(\exists q)(0 \leq q < a \wedge B[p]=A[q]+h(p) \wedge A[q]>0)$

Automating Program Reasoning

```
cnt=0, fib1=1, fib2=0;
```

```
while (cnt<n) do
```

```
t=fib1; fib1=fib1+fib2; fib2=t; cnt++;
```

```
end do
```

h

```
a=0, b=0, c=0;
```

```
while (a<n) do
```

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if A[a]>0 then B[b]=A[a]+h(b); b=b+1;
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  else C[c]=A[a]; c=c+1;
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end do
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Automating Program Reasoning

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```
end do
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Program property:

$$\text{fib1}^4 + \text{fib2}^4 + 2 * \text{fib1} * \text{fib2}^3 - 2 \text{fib1}^3 * \text{fib2} - \text{fib1}^2 * \text{fib2}^2 - 1 = 0$$

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a=0, b=0, c=0;
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while (a<n) do
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if A[a]>0 then B[b]=A[a]+h(b); b=b+1;
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end do
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Automating Program Reasoning

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$(\forall p)(0 \leq p < b \Rightarrow$

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Computer
Algebra

First-Order
Theorem Proving

My Research Group

Automated Program Reasoning - APro

Loop Analysis

Computer
Algebra

First-Order
Theorem Proving

My Research Group

Automated Program Reasoning - APRe

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Program Analysis

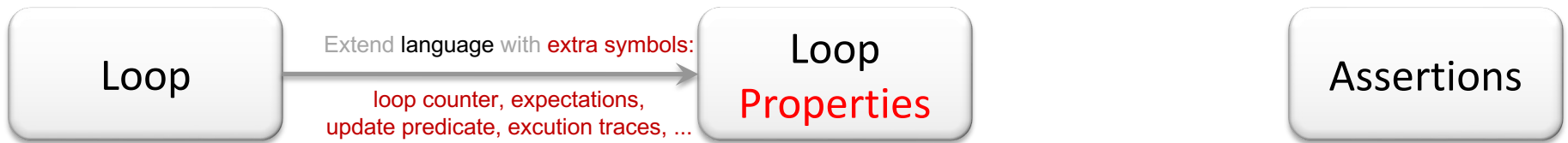
Our Recipe in Automated Program Reasoning

Our Recipe in Automated Program Reasoning

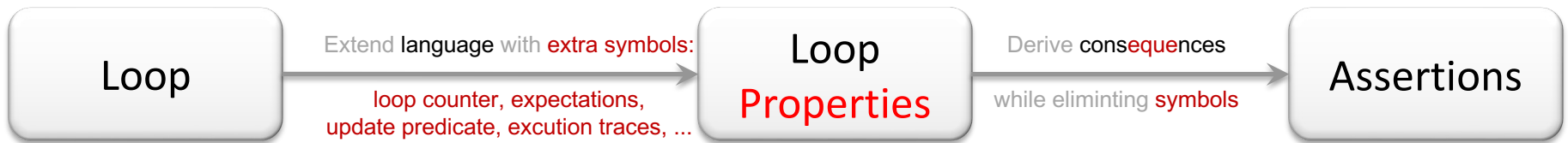
Loop

Assertions

Our Recipe in Automated Program Reasoning



Our Recipe in Automated Program Reasoning



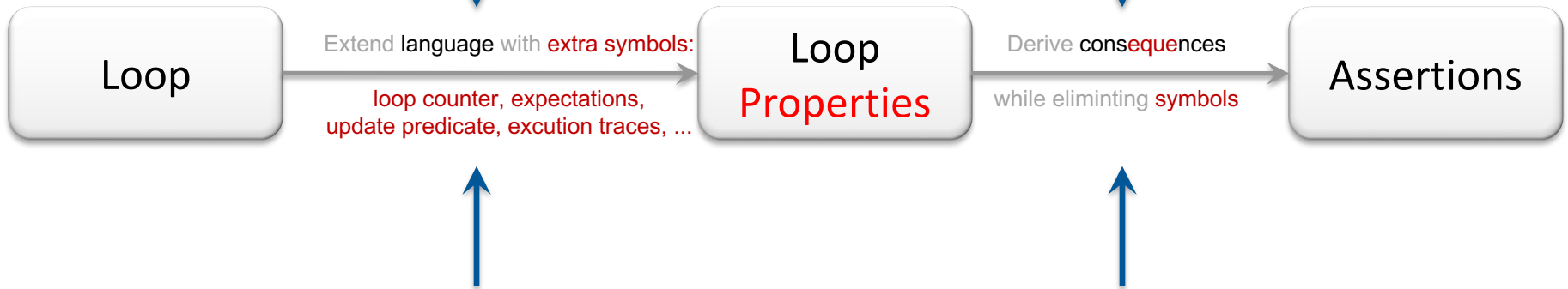
Our Recipe in Automated Program Reasoning

Algebraic recurrences

Statistical moments

Gröbner basis

Quantifier elimination



Loop

Extend language with extra symbols:
loop counter, expectations,
update predicate, excution traces, ...

Loop
Properties

Derive consequences
while eliminting symbols

Assertions

Extensionality axioms

Trace lemmas

Saturation

Induction

Symbolic Computation in Automated Program Reasoning

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Symbolic Computation in Automated Program Reasoning

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Loop

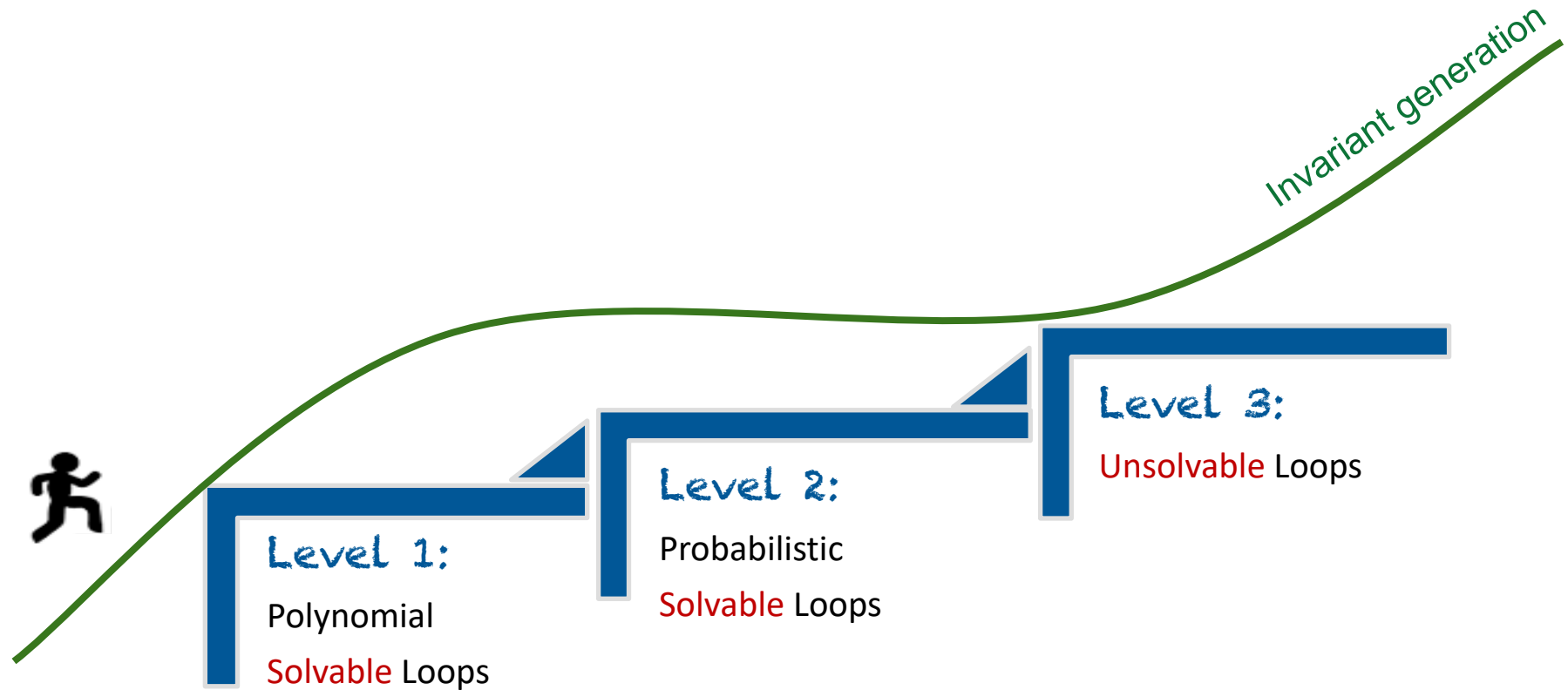
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Derive **consequences**
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Assertions

Symbolic Computation in Automated Program Reasoning



Symbolic Computation in Automated Program Reasoning

Level 1: Polynomial Solvable Loops

```
x:=1; y:=0;
```

```
while ... do x:=2*x; y:= $\frac{1}{2}$ *y+1 end do
```

Symbolic Computation in Automated Program Reasoning

Level 1: Polynomial Solvable Loops

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x:=1; y:=0;
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while ... do x:=2*x; y:= $\frac{1}{2}$ *y+1 end do
```

$n \geq 0$

1. Express state from $(n+1)^{\text{th}}$ iteration in terms of the n^{th} iteration \rightarrow algebraic recurrences of loop variables

$$\begin{cases} x(n+1) = 2 * x(n) \\ y(n+1) = \frac{1}{2} * y(n) + 1 \end{cases}$$

Symbolic Computation in Automated Program Reasoning

Level 1: Polynomial Solvable Loops

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2. Solve recurrences \rightarrow closed forms of loop variables

$$\begin{cases} x(n) = 2^n * x(0) \\ y(n) = \frac{1}{2^n} * y(0) - \frac{2}{2^n} + 2 \end{cases}$$

Symbolic Computation in Automated Program Reasoning

Level 1: Polynomial Solvable Loops

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x:=1; y:=0;  
while ... do x:=2*x; y:= $\frac{1}{2}$ *y+1 end do
```

$n \geq 0$, $a=2^n$, $b=2^{-n}$

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3. Derive algebraic dependencies among exponentials in n

$$\begin{cases} x(n) = a * x(0). \\ y(n) = b * y(0) - 2 * b + 2 \\ 0 = a * b - 1 = 2^n * \frac{1}{2^n} - 1 \end{cases}$$

Symbolic Computation in Automated Program Reasoning

Level 1: Polynomial Solvable Loops

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while ... do x:=2*x; y:= $\frac{1}{2}$ *y+1 end do
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4. Eliminate expressions in n \leftarrow Gröbner basis computation

$$x * y - 2 * x + 2 = 0$$

Symbolic Computation in Automated Program Reasoning

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4. Eliminate expressions in n \leftarrow Gröbner basis computation

$$x * y - 2 * x + 2 = 0$$

\rightarrow Finite basis of polynomial invariant ideal

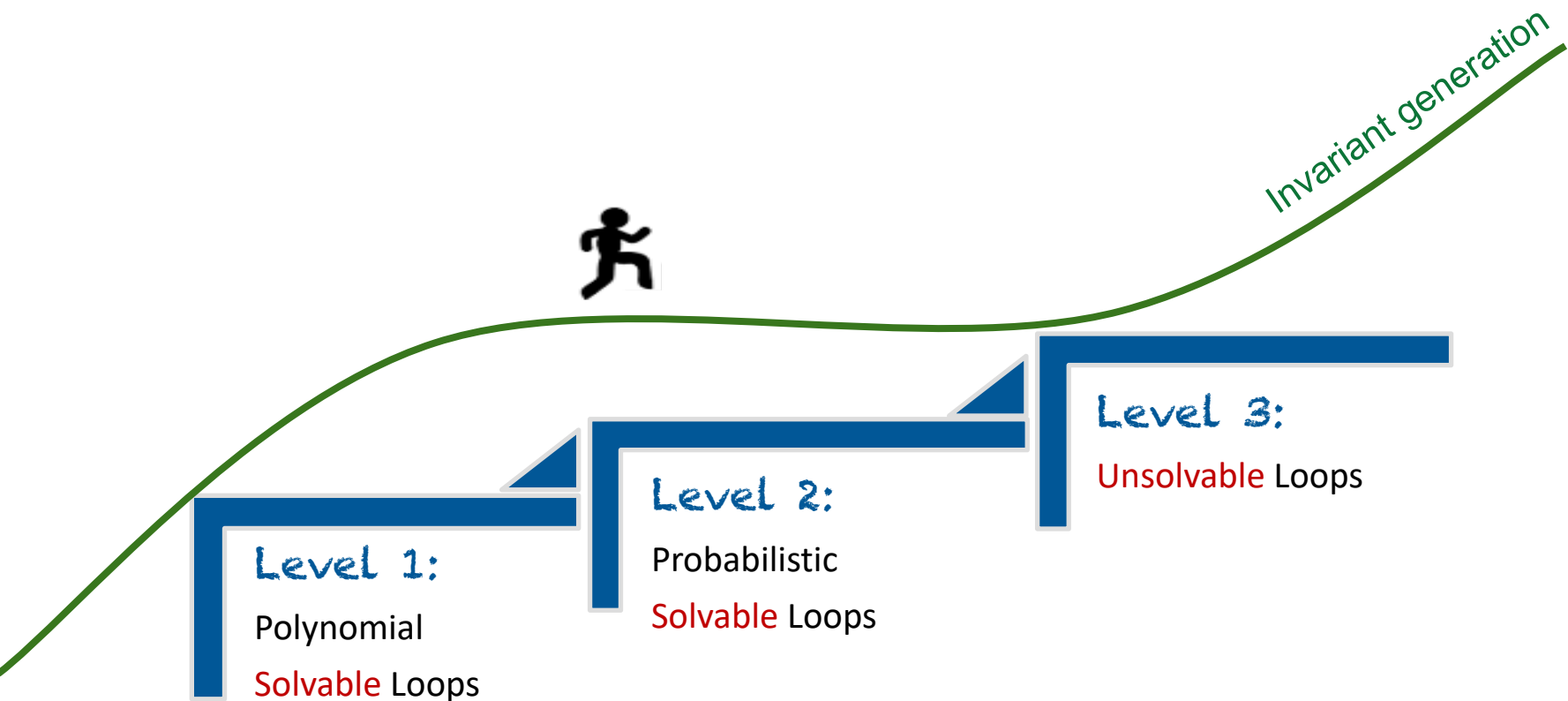
Symbolic Computation in Automated Program Reasoning

Level 1: Polynomial Solvable Loops

joint work w A. Humenberger, M. Jaroschek, A. Varonka (ISSAC17,VMCAI18,RAMICS23)

- Loops with polynomial assignments and nested conditionals
 - Structural constraints on assignments with polynomial rhs
 - ← C-finite recurrences of loop variables
 - Tests are ignored → non-deterministic programs
- Automation via symbolic summation and Gröbner basis computation
 - ALIGATOR tool <https://ahumenberger.github.io/aligator/>
- Further applications: loop termination, synthesis, deductive verification

Symbolic Computation in Automated Program Reasoning



Symbolic Computation in Automated Program Reasoning

Level 2: Probabilistic Solvable Loops

```
x:=0;  
while ... do  
  x:=x-1 [1/2] x+1;  
end do
```

- What is the **behaviour of a probabilistic loop**?

Symbolic Computation in Automated Program Reasoning

Level 2: Probabilistic Solvable Loops

```
x:=0;  
while ... do  
  x:=x-1 [1/2] x+1;  
end do
```

- What is the **behaviour of a probabilistic loop**?
- What is the **expected value** of a loop variable, e.g. x?

Symbolic Computation in Automated Program Reasoning

Level 2: Probabilistic Solvable Loops

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x:=0;  
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➤ $E[x]=0$

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```

➤ $E[x]=0$

```
x:=0;  
while ... do  
  x:=2*x-1 [1/2] 2*x+1;  
end do
```

- What is the **behaviour of a probabilistic loop**?
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Symbolic Computation in Automated Program Reasoning

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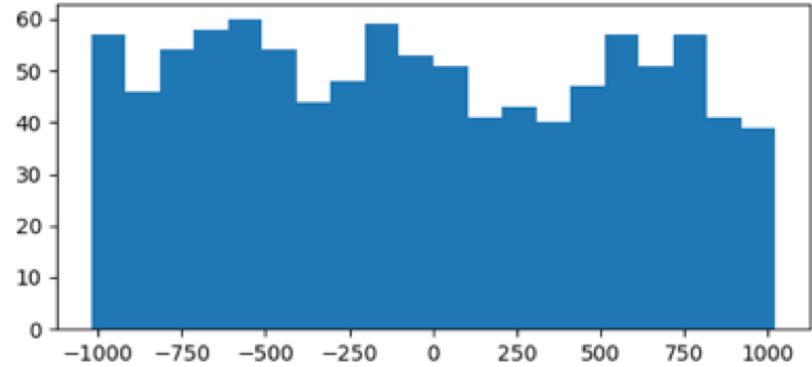
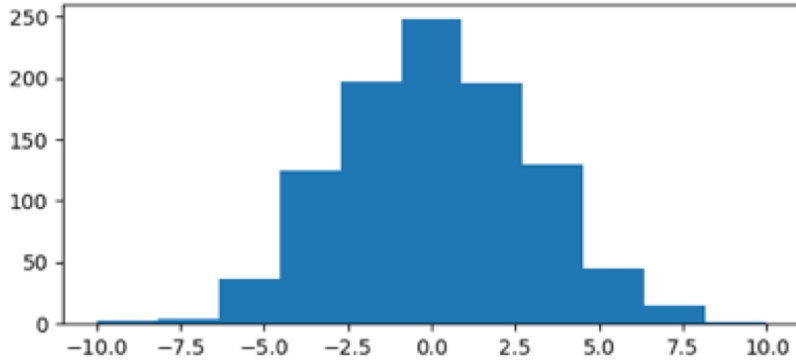
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x:=0;  
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➤ $E[x]=0$

- What is the **behaviour of a probabilistic loop**?
- What is the **expected value** of a loop variable, e.g. x ?
- In both programs above, the **expected value** of x is the same. Yet, the programs are **not the same!**

Symbolic Computation in Automated Program Reasoning

Level 2: Probabilistic Solvable Loops



- What is the **behaviour of a probabilistic loop**?
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Symbolic Computation in Automated Program Reasoning

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➤ $E[x]=0$

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- What is the **behaviour of a probabilistic loop**?
- Can we **characterize/recover the value distribution** of loop variables?

Symbolic Computation in Automated Program Reasoning

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- What is the **behaviour of a probabilistic loop**?
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Reason about higher-order statistical moments of variables!

Symbolic Computation in Automated Program Reasoning

Level 2: Probabilistic Solvable Loops

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x:=0;  
while ... do  
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```

➤ $E[x(n)]=0, \text{Var}[x(n)]=\frac{4^n}{3}$

```
x:=0;  
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➤ $E[x]=0, \text{Var}[x(n)]=n$

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Symbolic Computation in Automated Program Reasoning

Level 2: Probabilistic Solvable Loops

joint work w E. Bartocci, M. Stankovic, M. Moosbrugger (ATVA19, TACAS20, OOPSLA22)

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DEFINITION 14 (MOMENT-COMPUTABILITY). A probabilistic loop \mathcal{P} is moment-computable if a closed-form of $\mathbb{E}(x_n^k)$ exists and is computable for all $x \in \text{Vars}(\mathcal{P})$ and $k \in \mathbb{N}$.

Symbolic Computation in Automated Program Reasoning

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DEFINITION 14 (MOMENT-COMPUTABILITY). A probabilistic loop \mathcal{P} is moment-computable if a closed-form of $\mathbb{E}(x_n^k)$ exists and is computable for all $x \in \text{Vars}(\mathcal{P})$ and $k \in \mathbb{N}$.

Theorem 6 (Moment-Computability). A probabilistic loop \mathcal{P} is moment-computable if (1) none of its non-finite variables depends on itself polynomially, and (2) if the variables in all if-conditions are finite.

Symbolic Computation in Automated Program Reasoning

Level 2: Probabilistic Solvable Loops

joint work w E. Bartocci, M. Stankovic, M. Moosbrugger (ATVA19, TACAS20, OOPSLA22)

```
real x := -1, y := 1;  
real s := 0, f := 0, d;
```

A

```
while (true){
```

```
  f := 1 [3/4] 0;
```

```
  x := x + f * rand(1-d, 1+d);
```

```
  y := y + f * rand(2-2d, 2+2d);
```

```
  s := x + y;
```

```
}
```

$$E[s_n] = \frac{9}{4}n$$

$$\text{Var}[s_n] = \frac{20d^2 + 27}{16}n$$

```
real x := rand (-9, 7), y := rand (-7, 9);  
real s := 0, f := 0;
```

B

```
while (true){
```

```
  f := 1 [3/4] 0;
```

```
  x := x + f * rand(-3,5);
```

```
  y := y + f * rand(-6,10);
```

```
  s := x + y;
```

```
}
```

$$E[s_n] = \frac{9}{4}n$$

$$\text{Var}[s_n] = \frac{347}{16}n + \frac{128}{3}$$

- Parametrized distributions
 - Random polynomial assignments → C-finite recurrences over moments
 - Finite-valued multi-path conditions
- Higher-order moments of loop variables are computable.

Symbolic Computation in Automated Program Reasoning

Level 2: Probabilistic Solvable Loops

joint work w E. Bartocci, M. Stankovic, M. Moosbrugger (ATVA19, TACAS20, OOPSLA22)

```
f, x, y, s = 0, -1, 1, 0
while (true):
    f = 1 [3/4] 0
    x = x + f*rand(1-d, 1+d)
    y = y + f*rand(2-2d, 2+2d)
    s = x + y
```

Goal: Higher-order Moments,
e.g. s^2

Symbolic Computation in Automated Program Reasoning

Level 2: Probabilistic Solvable Loops

joint work w E. Bartocci, M. Stankovic, M. Moosbrugger (ATVA19, TACAS20, OOPSLA22)

- **Expected values** of monomials

- sufficient to understand **E-variables**

$E[X], E[X^2], E[XY], \dots$

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```

Probabilistic updates:

$$x_i = a_i x_i + P_i(x_1, \dots, x_{i-1}) [p_i] b_i x_i + Q_i(x_1, \dots, x_{i-1})$$

Symbolic Computation in Automated Program Reasoning

Level 2: Probabilistic Solvable Loops

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Stochastic recurrences over E-variables:

$$E[x_i(n+1)] = p_i \cdot E[a_i x_i(n) + P_i(x_1(n), \dots, x_{i-1}(n))] \\ + (1 - p_i) \cdot E[b_i x_i(n) + Q_i(x_1(n), \dots, x_{i-1}(n))].$$

Symbolic Computation in Automated Program Reasoning

Level 2: Probabilistic Solvable Loops

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```

- Computing with **E-variables**

$$E[c] = c$$

$$E[X + cY] = E[X] + cE[Y]$$

$$E[X \cdot Y] \neq E[X] \cdot E[Y] \text{ unless } X, Y \text{ are independent}$$

Stochastic recurrences over E-variables:

$$E[x_i(n+1)] = p_i \cdot E[a_i x_i(n) + P_i(x_1(n), \dots, x_{i-1}(n))] \\ + (1 - p_i) \cdot E[b_i x_i(n) + Q_i(x_1(n), \dots, x_{i-1}(n))].$$

Symbolic Computation in Automated Program Reasoning

Level 2: Probabilistic Solvable Loops

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- **Expected values** of monomials
- **necessary** to understand **E-variables**

$$E[X], E[X^2], E[XY], \dots$$

```
f, x, y, s = 0, -1, 1, 0
while (true):
    f = 1 [3/4] 0
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$$E[c] = c$$

$$E[X + cY] = E[X] + cE[Y]$$

$$E[X \cdot Y] \neq E[X] \cdot E[Y] \text{ unless } X, Y \text{ are independent}$$

Expected values of monomials cannot be simplified

Stochastic recurrences over E-variables:

$$E[x_i(n+1)] = p_i \cdot E[a_i x_i(n) + P_i(x_1(n), \dots, x_{i-1}(n))] \\ + (1 - p_i) \cdot E[b_i x_i(n) + Q_i(x_1(n), \dots, x_{i-1}(n))].$$

Symbolic Computation in Automated Program Reasoning

Level 2: Probabilistic Solvable Loops

joint work w E. Bartocci, M. Stankovic, M. Moosbrugger (ATVA19, TACAS20, OOPSLA22)

- $x(n+1) = 2 * x(n)$



Level 1

- $x(n+1) = x(n) + \text{unif}(0,1)$

Symbolic Computation in Automated Program Reasoning

Level 2: Probabilistic Solvable Loops

joint work w E. Bartocci, M. Stankovic, M. Moosbrugger (ATVA19, TACAS20, OOPSLA22)

- $x(n+1) = 2 * x(n)$



Level 1

- $x(n+1) = x(n) + \text{unif}(0,1)$



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Level 2: Probabilistic Solvable Loops

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- $x(n+1) = 2 \cdot x(n)$



Level 1

- $x(n+1) = x(n) + \text{unif}(0,1)$



—use **moments**: treat each moment as a separate **E-variable**

- $E[x(n+1)] = E[x(n)] + 1/2$

$$E[x^2(n+1)] = E[x^2(n)] + \dots$$

$$E[x^3(n+1)] = \dots$$

...

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Level 1

- $x(n+1) = x(n) + \text{unif}(0,1)$



—use **moments**: treat each moment as a separate **E-variable**

- $E[x(n+1)] = E[x(n)] + 1/2$
 $E[x^2(n+1)] = E[x^2(n)] + \dots$
 $E[x^3(n+1)] = \dots$
 \dots

—solve **stochastic recurrences**: closed forms over **E-variables**

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f, x, y, s = 0, -1, 1, 0
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    f = 1 [3/4] 0
    x = x + f*rand(1-d, 1+d)
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```

goal: {s^2}

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$\{s^2\}$

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goal: {s^2}
```

$$\{s^2\} \rightarrow s^2$$

Symbolic Computation in Automated Program Reasoning

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goal: {s^2}
```

$$\{s^2\} \rightarrow s^2 \rightarrow E[s^2(n+1)] = E[(x(n+1) + y(n+1))^2]$$

Symbolic Computation in Automated Program Reasoning

Level 2: Probabilistic Solvable Loops

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goal: {s^2}
```

$$\{s^2\} \rightarrow s^2 \rightarrow E[s^2(n+1)] = E[(x(n+1) + y(n+1))^2]$$
$$\rightarrow E[s^2(n+1)] = E[x^2(n+1)] + 2E[xy(n+1)] + E[y^2(n+1)]$$

Symbolic Computation in Automated Program Reasoning

Level 2: Probabilistic Solvable Loops

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```

$$\{s^2\} \rightarrow s^2 \rightarrow E[s^2(n+1)] = E[(x(n+1) + y(n+1))^2]$$
$$\rightarrow E[s^2(n+1)] = E[x^2(n+1)] + 2E[xy(n+1)] + E[y^2(n+1)]$$
$$\rightarrow S = \{x^2, xy, y^2\}$$

Symbolic Computation in Automated Program Reasoning

Level 2: Probabilistic Solvable Loops

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    x = x + f*rand(1-d, 1+d)
    y = y + f*rand(2-2d, 2+2d)
    s = x + y

goal: {s^2}
```

→ x^2

Symbolic Computation in Automated Program Reasoning

Level 2: Probabilistic Solvable Loops

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    y = y + f*rand(2-2d, 2+2d)
    s = x + y

goal: {s^2}
```

$$\rightarrow x^2 \rightarrow E[x^2(n+1)] = E[(x(n) + f(n+1) \cdot \text{rand}(1-d, 1+d))^2]$$

Symbolic Computation in Automated Program Reasoning

Level 2: Probabilistic Solvable Loops

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$$\rightarrow x^2 \rightarrow E[x^2(n+1)] = E[(x(n) + f(n+1) \cdot \text{rand}(1-d, 1+d))^2]$$

$$\rightarrow E[x^2(n+1)] = E[x^2(n)] + 2E[x(n)]E[f(n+1)] + E[(f(n+1))^2](1 + \frac{d^2}{3})$$

Symbolic Computation in Automated Program Reasoning

Level 2: Probabilistic Solvable Loops

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goal: {s^2}
```

$$\rightarrow x^2 \rightarrow E[x^2(n+1)] = E[(x(n) + f(n+1) \cdot \text{rand}(1-d, 1+d))^2]$$

$$\rightarrow E[x^2(n+1)] = E[x^2(n)] + 2E[x(n)]E[f(n+1)] + E[(f(n+1))^2](1 + \frac{d^2}{3})$$

$$\rightarrow S = \{xy, y^2, x, f, f^2\}$$

Symbolic Computation in Automated Program Reasoning

Level 2: Probabilistic Solvable Loops

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2. While S is not empty:
3. Pick an E-variable from S
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    f = 1 [3/4] 0
    x = x + f*rand(1-d, 1+d)
    y = y + f*rand(2-2d, 2+2d)
    s = x + y

goal: {s^2}
```

$$\rightarrow x^2 \rightarrow E[x^2(n+1)] = E[(x(n) + f(n+1) \cdot \text{rand}(1-d, 1+d))^2]$$

$$\rightarrow E[x^2(n+1)] = E[x^2(n)] + 2E[x(n)]E[f(n+1)] + E[(f(n+1))^2](1 + \frac{d^2}{3})$$

$$\rightarrow S = \{xy, y^2, x, f, f^2\}$$

...and so on ...

Symbolic Computation in Automated Program Reasoning

Level 2: Probabilistic Solvable Loops

joint work w E. Bartocci, M. Stankovic, M. Moosbrugger (ATVA19, TACAS20, OOPSLA22)

$$E[f(n + 1)] = \frac{3}{4}$$

$$E[x(n + 1)] = E[x(n)] + \frac{3}{4}E[f(n + 1)]$$

$$E[y(n + 1)] = E[y(n)] + \frac{3}{2}E[f(n + 1)]$$

$$E[s(n + 1)] = E[x(n + 1)] + E[y(n + 1)]$$

...

$$E[s(n + 1)^2] = E[x(n + 1)^2] + 2 \cdot E[(xy)(n + 1)] + E[y(n + 1)^2]$$

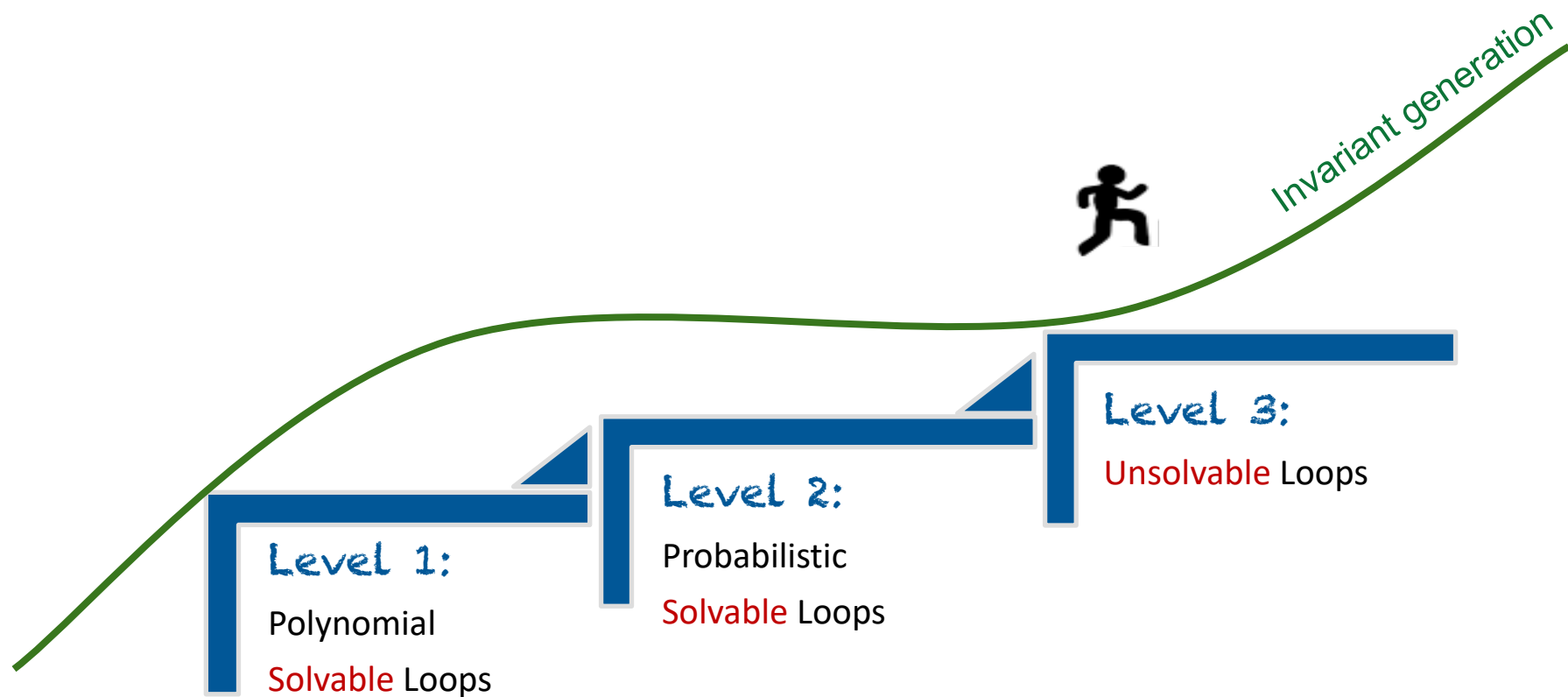
Symbolic Computation in Automated Program Reasoning

Level 2: Probabilistic Solvable Loops

joint work w A. Humenberger, M. Jaroschek, A. Varonka (ISSAC17,VMCAI18,RAMICS18)

- Level 1 + Random polynomial assignments + finite-valued conditions
 - C-finite recurrences of E-variables
 - Tests are finite-valued
- Automation via symbolic summation and moment-based computation
 - POLAR tool <https://github.com/probing-lab/polar>
- Further applications: probabilistic termination, sensitivity, probabilistic inferences

Symbolic Computation in Automated Program Reasoning



Symbolic Computation in Automated Program Reasoning

Level 3: Unsolvable Loops

```
a:=-2; b:=3; y:=0;
```

```
while ... do
```

```
  a:=2*a+b2;
```

```
  b:=2*b-b2;
```

```
  y:= $\frac{1}{2}$ *y+1
```

```
end do
```


Symbolic Computation in Automated Program Reasoning

Level 3: Unsolvable Loops

```
a:=-2; b:=3; y:=0;  
while ... do  
  a:=2*a+b2;  
  b:=2*b-b2;  
  y:= $\frac{1}{2}$ *y+1  
end do
```

➤ Polynomial updates → non-C-finite recurrences

Level 1 and Level 2 ⊗

Symbolic Computation in Automated Program Reasoning

Level 3: Unsolvable Loops

```
a:=-2; b:=3; y:=0;  
while ... do  
  a:=2*a+b2;  
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  y:= $\frac{1}{2}$ *y+1  
end do
```

$n \geq 0$

- Polynomial updates → non-C-finite recurrences Level 1 and Level 2 ⊗
- Yet, $x(n)=a(n)+b(n)$ satisfy a C-finite recurrence: $a(n+1)+b(n+1)=2*(a(n)+b(n))$

Symbolic Computation in Automated Program Reasoning

Level 3: Unsolvable Loops

→

Level 1: Solvable Loops

```
a:=-2; b:=3; y:=0;
while ... do
  a:=2*a+b2;
  b:=2*b-b2;
  y:= $\frac{1}{2}$ *y+1
end do
```

$n \geq 0$

```
x:=1; y:=0;
while ... do
  x:=2*x;
  y:= $\frac{1}{2}$ *y+1
end do
```

- Polynomial updates → non-C-finite recurrences Level 1 and Level 2 ⊗
- Yet, $x(n)=a(n)+b(n)$ satisfy a C-finite recurrence: $a(n+1)+b(n+1)=2*(a(n)+b(n))$
- Unsolvable loop over a, b, y → Solvable loop over x, y Level 1 ✓

Symbolic Computation in Automated Program Reasoning

Level 3: Unsolvable Loops

→

Level 1: Solvable Loops

```
a:=-2; b:=3; y:=0;
while ... do
  a:=2*a+b2;
  b:=2*b-b2;
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end do
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$n \geq 0$

```
x:=1; y:=0;
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➤ Polynomial updates → non-C-finite recurrences Level 1 and Level 2 ⊗

➤ Yet, $x(n)=a(n)+b(n)$ satisfy a C-finite recurrence: $a(n+1)+b(n+1)=2*(a(n)+b(n))$

➤ Unsolvable loop over a, b, y → Solvable loop over x, y Level 1 ✓

Invariant: $(a+b)*y-2*(a+b)+2=0$ ← $x*y-2*x+2=0$

Symbolic Computation in Automated Program Reasoning

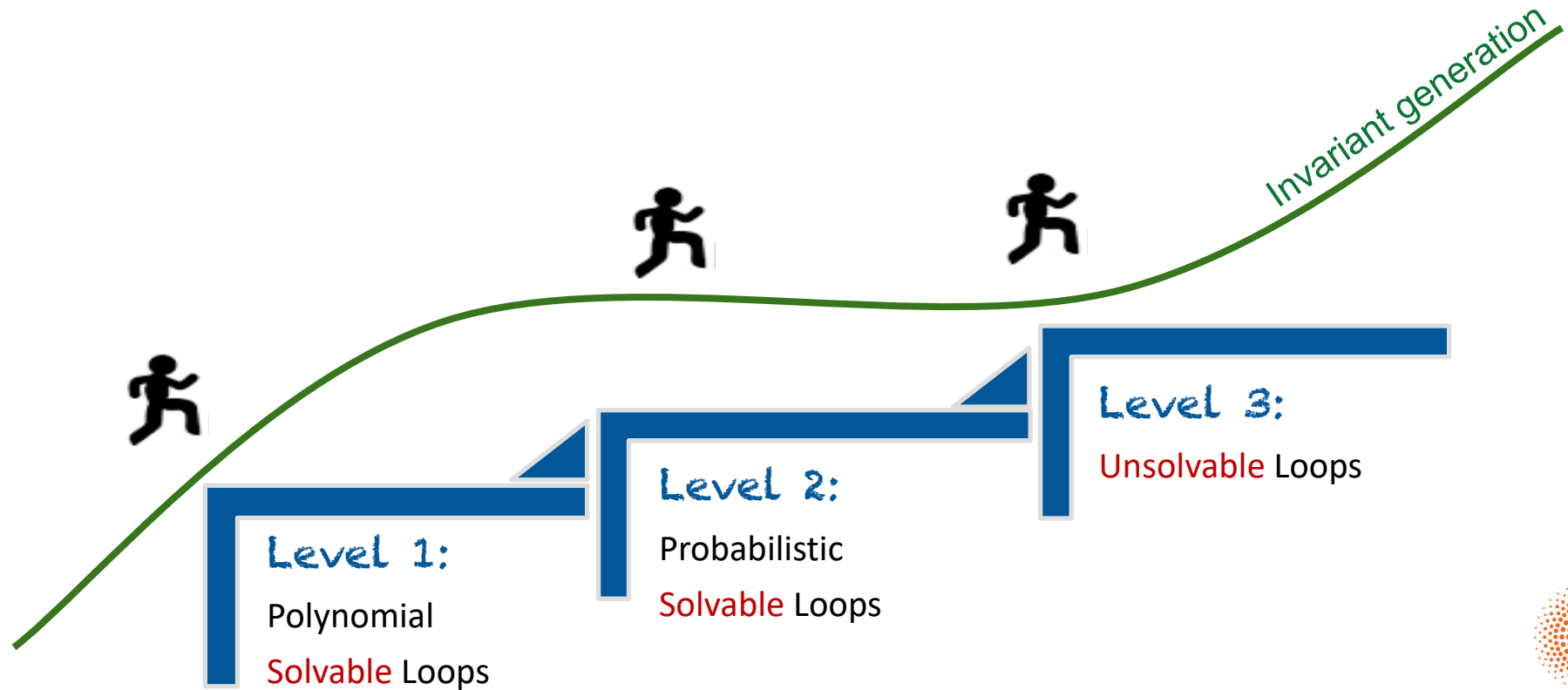
Level 3: **Unsolvable** Loops

→ Level 1 or Level 2: **Solvable** Loops

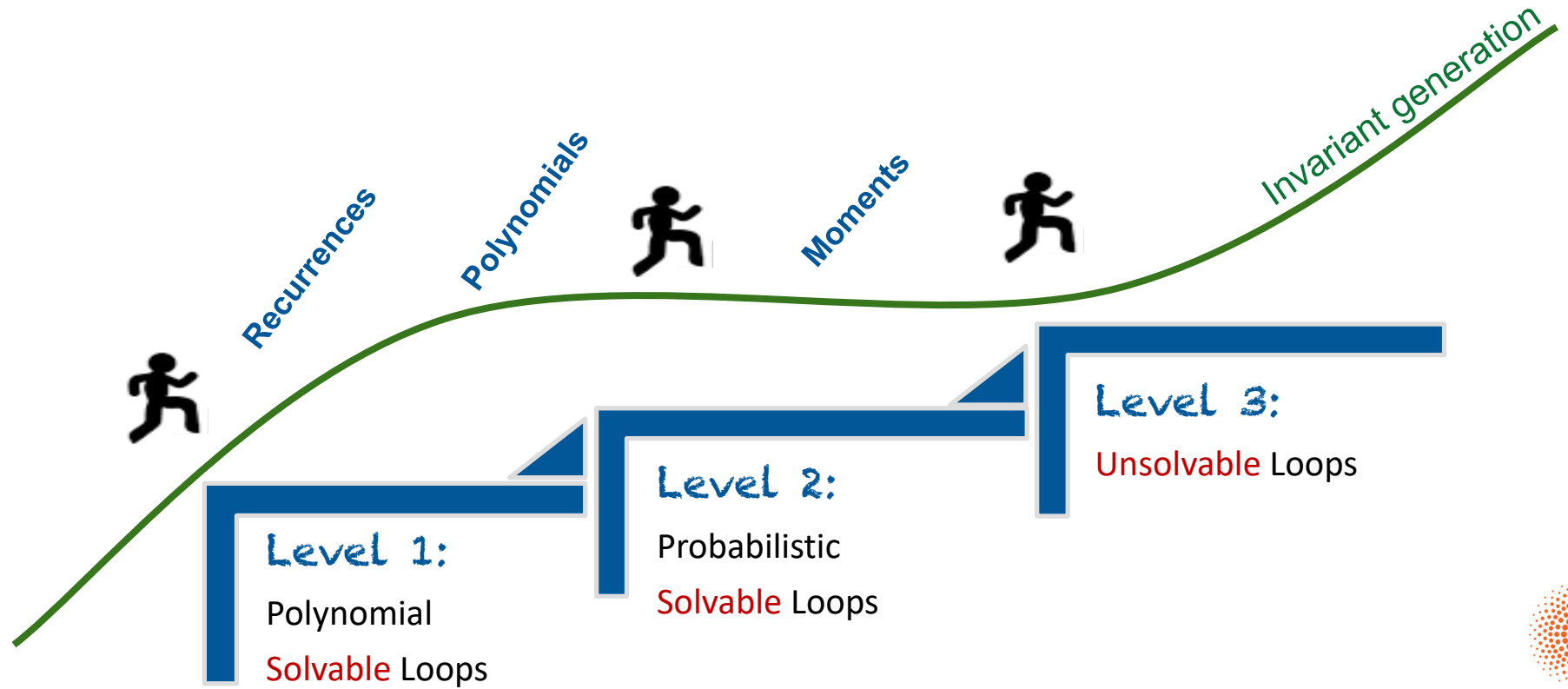
joint work w D. Amrollahi, E. Bartocci, G. Kenison, M. Stankovic, M. Moosbrugger (SAS22)

- Level 1 or Level 2 + **non-C-finite recurrences**
- Compute **polynomial relations P** over variables with non C-finite recurrences
- If P is C-finite expression, use P to **solve unsolvable loops**
- Automation via **variable dependency analysis** and **polynomial constraint solving**
 - **POLAR** tool <https://github.com/probing-lab/polar>

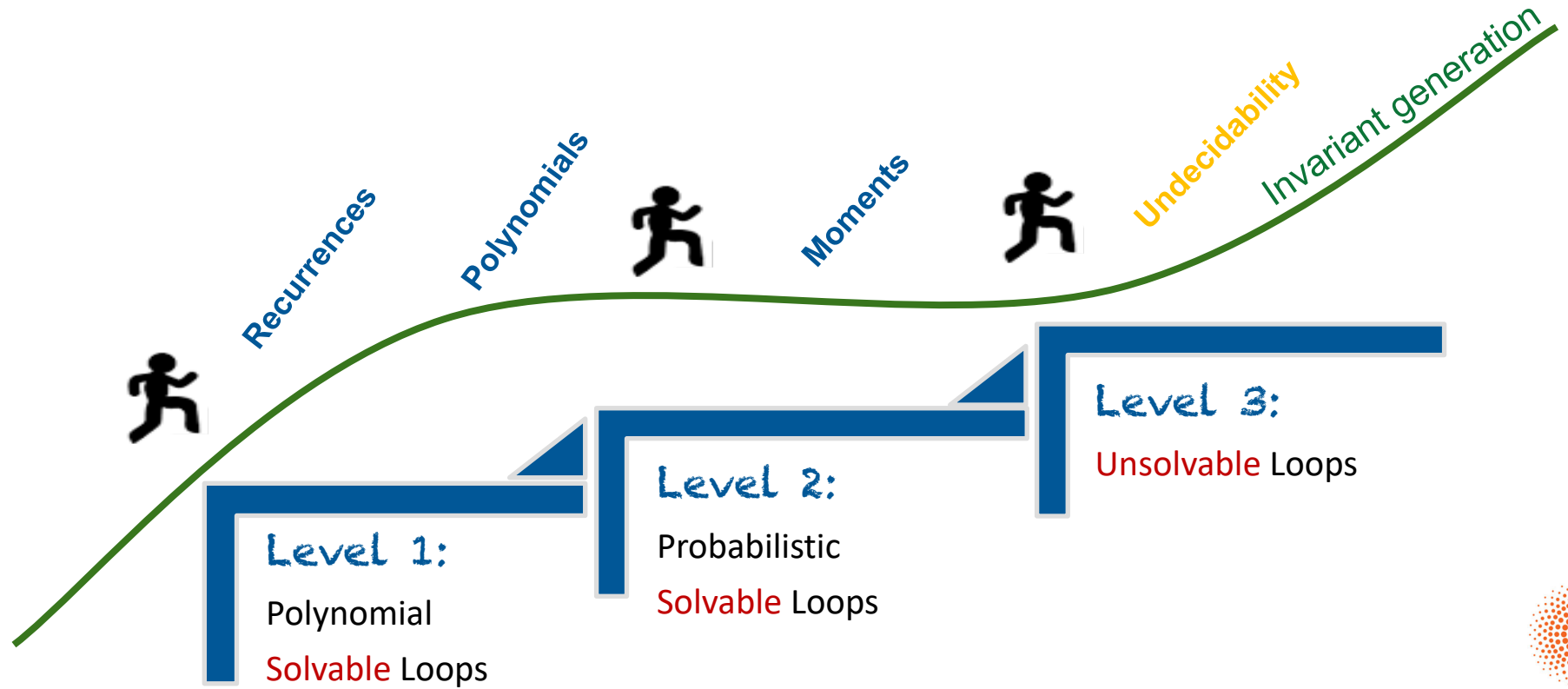
Symbolic Computation in Automated Program Reasoning



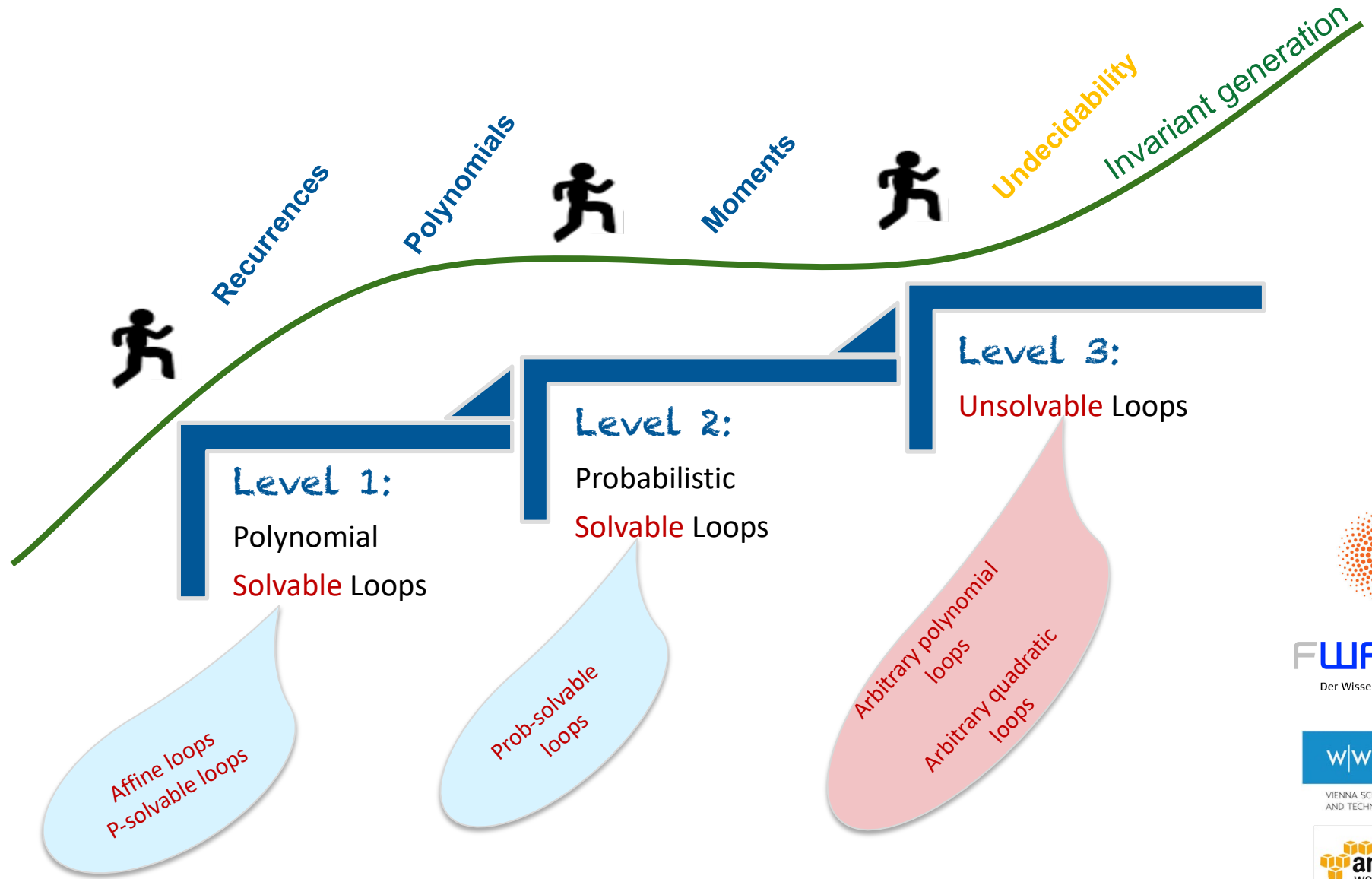
Symbolic Computation in Automated Program Reasoning



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